

Polyhedral Methods for Improving Parallel Update-in-Place

IMPACT'14

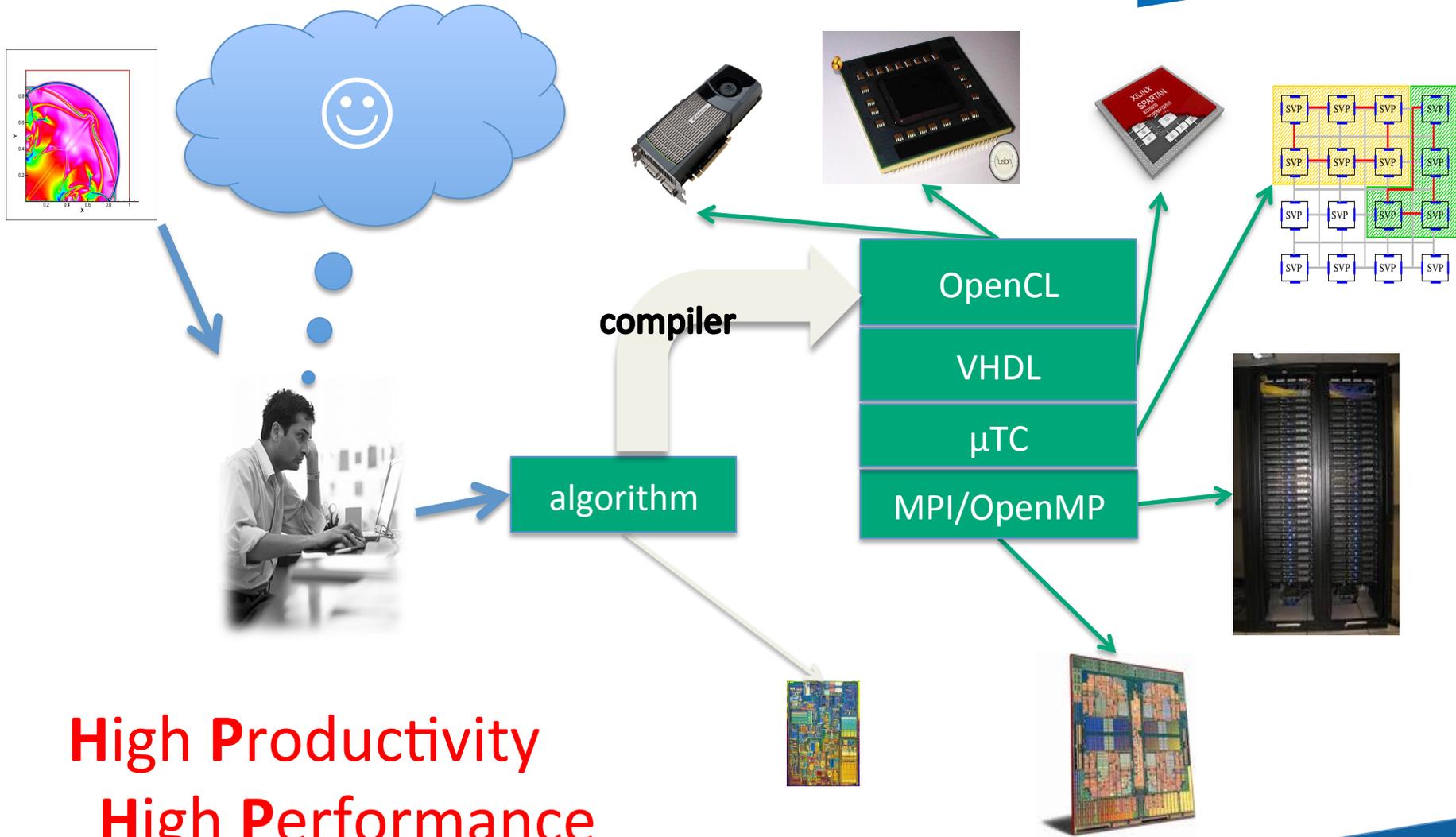
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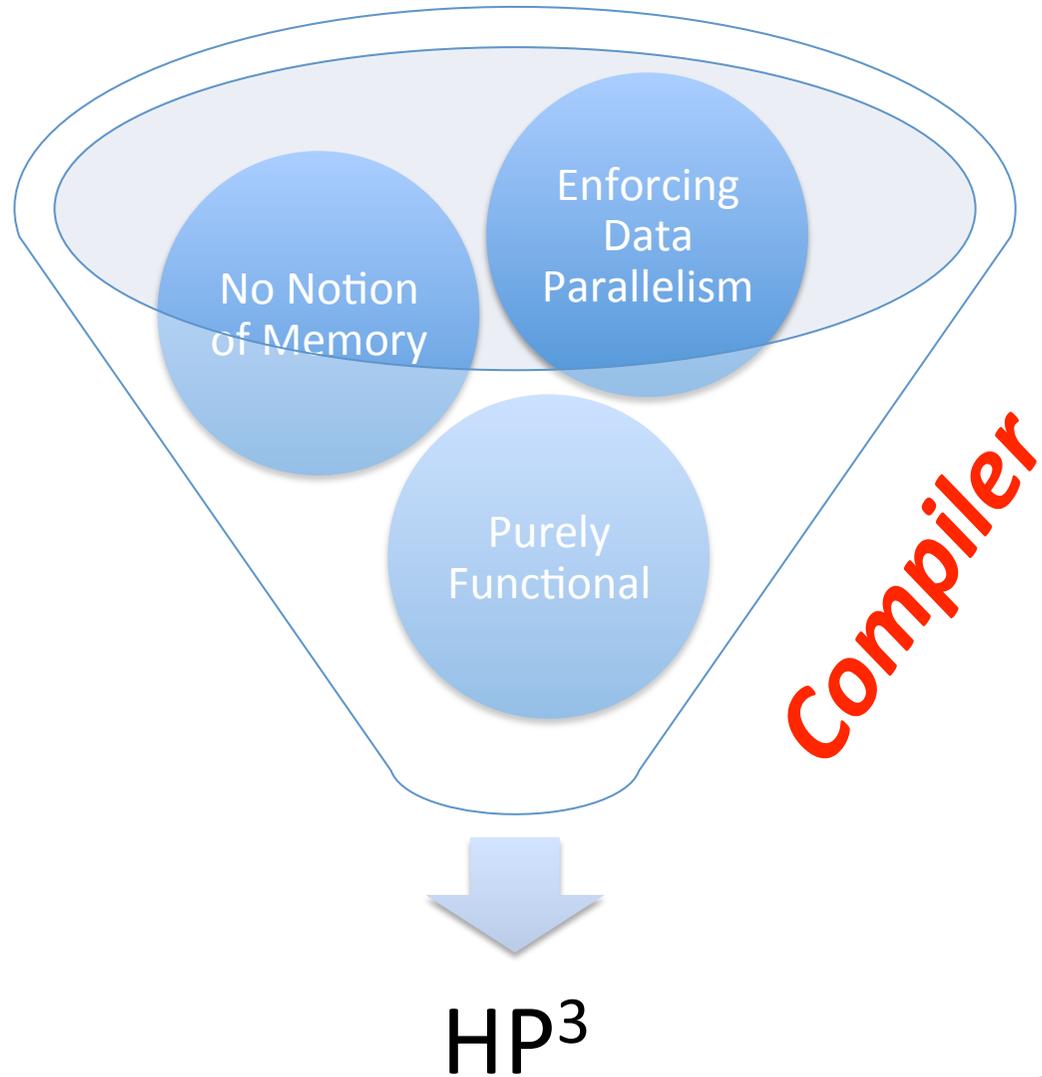
(University of Hertfordshire, Snake Island Research, Oxford e-
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HP³ Vision



High Productivity
High Performance
High Portability

“Magic” Ingredients



With-Loops in SaC

```
a = with {
    ( [1,1] <= iv < [3,3] ) : f(iv);
}: genarray( [5,8], def);
```

def	def	def	def	def	def	def	def
def	f([1,1])	f([1,2])	f([1,3])	def	def	def	def
def	f([2,1])	f([2,2])	f([2,3])	def	def	def	def
def	f([3,1])	f([3,2])	f([3,3])	def	def	def	def
def	def	def	def	def	def	def	def

Memory Management

```
a = with {  
    ( [1,1] <= iv < [3,3] ) : e(iv);  
} : genarray( [5,8], def);
```

allocate

Iterate & fill

```
b = ... a ..... a .....
```

```
c = ..... a .....
```

free

Memory Reuse Example

```
b = with {
    ( [1,1] <= iv < [3,3] ) : f(iv);
} : modarray( a);
```

a[iv]	a[iv]	a[iv]	a[iv]	a[iv]	a[iv]	a[iv]	a[iv]
a[iv]	f([1,1])	f([1,2])	f([1,3])	a[iv]	a[iv]	a[iv]	a[iv]
a[iv]	f([2,1])	f([2,2])	f([2,3])	a[iv]	a[iv]	a[iv]	a[iv]
a[iv]	f([3,1])	f([3,2])	f([3,3])	a[iv]	a[iv]	a[iv]	a[iv]
a[iv]	a[iv]	a[iv]	a[iv]	a[iv]	a[iv]	a[iv]	a[iv]

Can we reuse existing array a to represent new array b ?

- Smaller memory footprint
- Avoid copying overhead
- Depends on reference count

Memory Reuse ?

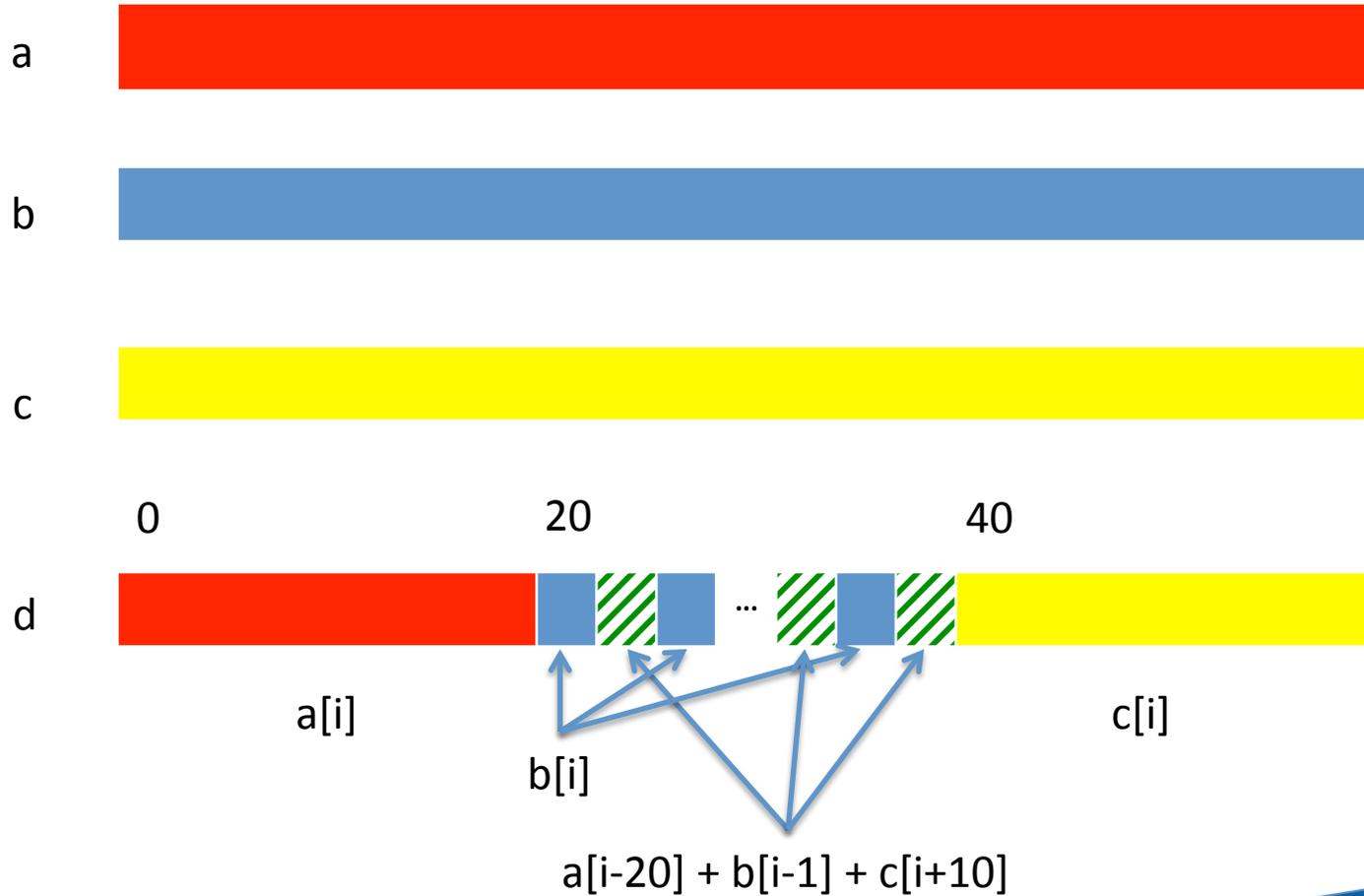
a[iv]	a[iv]	a[iv]	a[iv]	a[iv]	a[iv]	a[iv]	a[iv]
a[iv]	f([1,1])	f([1,2])	f([1,3])	a[iv]	a[iv]	a[iv]	a[iv]
a[iv]	f([2,1])	f([2,2])	f([2,3])	a[iv]	a[iv]	a[iv]	a[iv]
a[iv]	f([3,1])	f([3,2])	f([3,3])	a[iv]	a[iv]	a[iv]	a[iv]
a[iv]	a[iv]	a[iv]	a[iv]	a[iv]	a[iv]	a[iv]	a[iv]

- f(iv) = 42 ✓
- f(iv) = c[iv+1] ✓
- f(iv) = a[iv] ✓
- f(iv) = a[iv + [1,0]] ✗
- f(iv) = a[iv + [3,0]] ✓
- f(iv) = a[0,0] ✓

Can we recycle a ?

General Case

Which array (if any) can we recycle?



Observations



Polyhedral representations lend themselves!

Polyhedral tooling may be used!

Identifying *R*ead *D*ata *S*paces



```
4 foreach Array access  $\mathcal{A}[Iv]$  in  $\mathcal{Wl}$  do
5   Let  $\mathcal{I}$  be the iteration vector associated with this
   access;
6   Let CP be the set of control paths associated with
   this access;
7   if  $\mathcal{A}[Iv]$  is a non-in-place read access then
8     Affine,  $\mathcal{DS} \leftarrow \text{Get\_Data\_Space}(\mathcal{I}, Iv, \mathbf{CP},$ 
     SCiv);
9     if Affine = True then
10       $\mathbf{RDS}[\mathcal{A}] \leftarrow \mathbf{RDS}[\mathcal{A}] \cup \mathcal{DS};$ 
11    else
12      Terminate;
```

Identifying *C*opy *D*ata *S*paces



```
13  else if  $\mathcal{A}[Iv]$  is the global write access in a copy  
    assignment then  
14      Let  $\mathcal{A}'[Iv]$  be the in-place read of this copy  
        assignment;  
15       $Affine, DS \leftarrow \text{Get\_Data\_Space}(I, Iv, CP,$   
         $SCiv)$ ;  
16      if  $Affine = True$  then  
17         $\lfloor \text{CDS}[\mathcal{A}'] \leftarrow \text{CDS}[\mathcal{A}'] \cup DS;$   
18      else  
19         $\lfloor \text{Terminate};$   
20  else  
21     $\lfloor \text{Continue};$ 
```

Identifying *R*euse *C*andidates



```
22  $\mathbf{RC} \leftarrow \emptyset;$   
23 foreach  $\mathcal{A}$  with  $(\mathcal{A} \mapsto \mathcal{CDS}) \in \mathbf{CDS}$  do  
24   if  $\mathbf{RDS}[\mathcal{A}] \subseteq \mathbf{CDS}[\mathcal{A}]$  then  
25      $\mathbf{RC} \leftarrow \mathbf{RC} \cup \{\mathcal{A}\};$ 
```

Use the PolyLib for polyhedral operations:

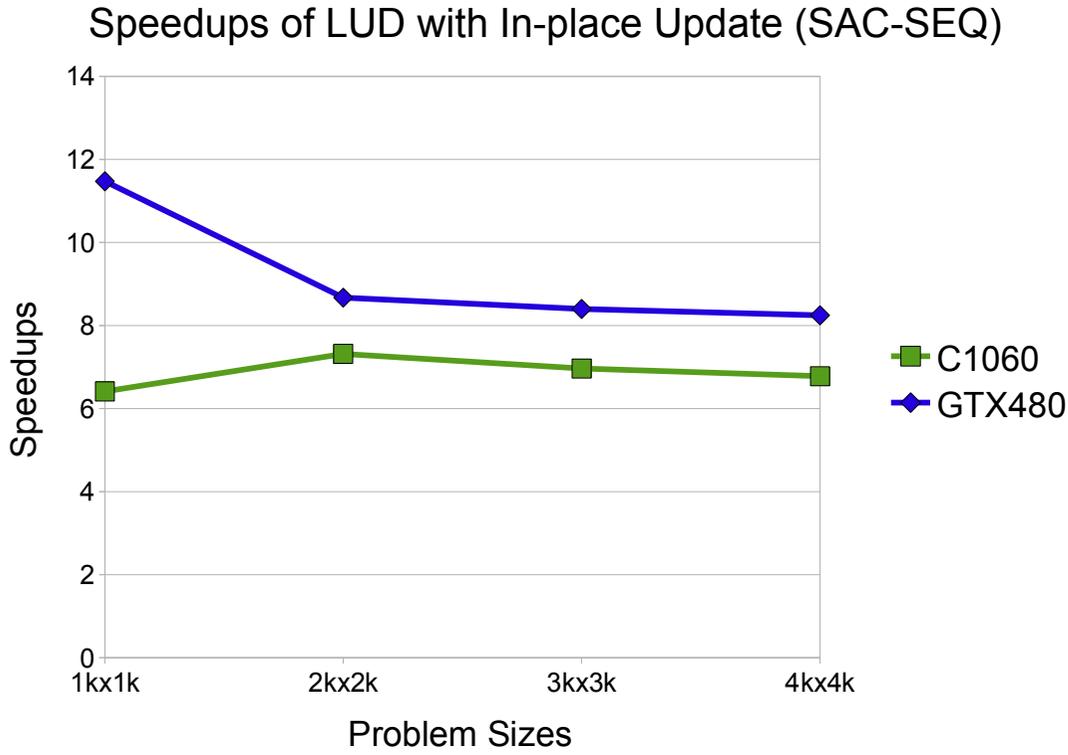
- Union of data spaces
- Inclusion relationship
- Image of polytope under affine transformation

Example: LU Decomposition



```
for( k = 0; k < N-1; k++) {  
2   A = with {  
      ( [k+1,k] <= [i , j] < [N,k+1])  
4           : A[i , j] / A[k , k];  
      } : modarray( A);  
6   A = with {  
      ( [k+1,k+1] <= [i , j] < [N,N])  
8           : A[i , j] - A[i , k] * A[k , j];  
      } : modarray( A);  
10 }
```

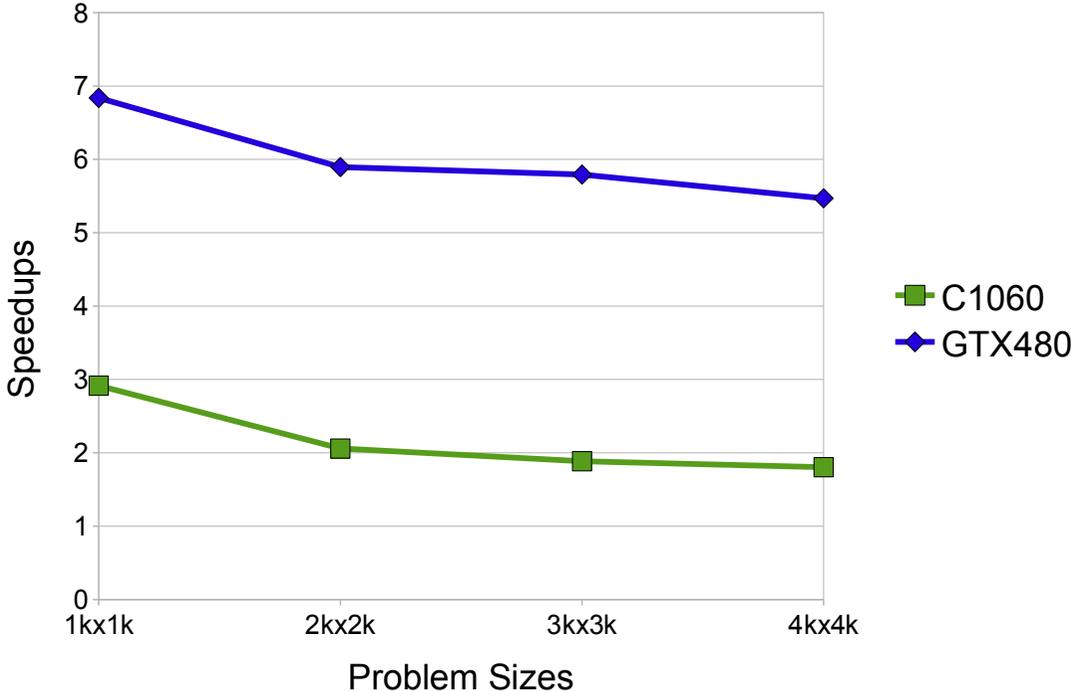
Runtime Impact (single core)



	-noPRA	-doPRA
1024 x 1024	0.28 GFLOPS	1.49 GFLOPS
2048 x 2048	0.19 GFLOPS	1.33 GFLOPS
4096 x 4096	0.19 GFLOPS	1.29 GFLOPS

Runtime Impact (GPU)

Speedups of LUD with In-place Update (SAC-CUDA)



	-noPRA	-doPRA
1024 x 1024	0.56 GFLOPS	10.21 GFLOPS
2048 x 2048	0.97 GFLOPS	11.22 GFLOPS
4096 x 4096	1.15 GFLOPS	11.39 GFLOPS

Memory Impact



	-noPRA	-doPRA
1024 x 1024	16 MB	9 MB
2048 x 2048	66 MB	34 MB
4096 x 4096	258 MB	138 MB

Conclusions

- Polyhedral model is an excellent means to formalise array accesses
- Even a dependency free setting benefits from the full power
- Reuse enables copy elimination
- Leads to essential speedups/ bridges the gap to classical for loops